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# Improving IEEE 802.15.4 for Low-latency Energy-efficient Industrial Applications

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# WSNs in Industrial Applications

## General requirements

- Standardization
- Coexistence, especially with WLAN
- Robust and reliable communications
- Energy-efficiency
- Security mechanism

## Why choose IEEE 802.15.4

- Industrial standard for LR-WPANs
- Coexistence with WLAN due to DSSS
- Duty-cycle adjustable
- GTS for real-time applications
- Provide security mechanism

# Study Case

## A typical industrial application

- A star network monitors industrial processes
  - e.g.: 20 sensor nodes and 1 gateway
- Nodes send alarm messages to the gateway
  - e.g.: a short alarm message with only 1 byte payload

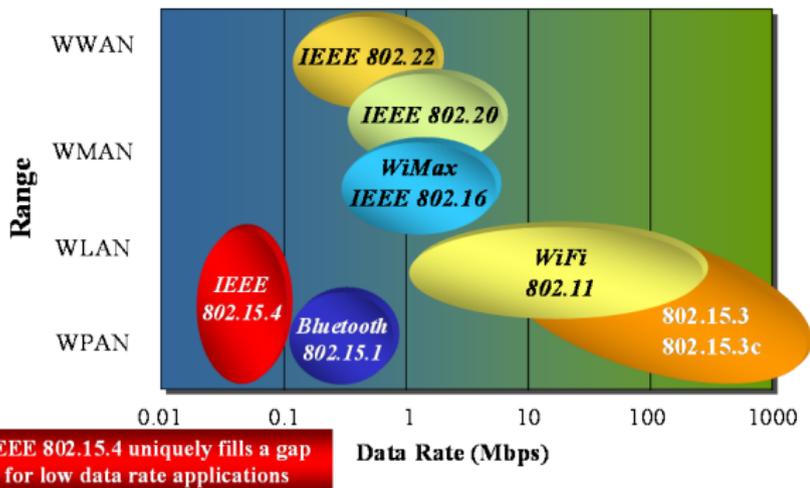
## Requirements

- Real-time: guaranteed latency upper boundary
  - e.g.:  $d_{GUA} < 10ms$
- Consideration of energy-efficiency
- Standardized hardwares

# Overview of IEEE 802.15.4

## Standard for Low-Rate WPANs

- low data rate  
(max. 250 kb/s)
- short distance  
(POS of 10m)
- ultra-low  
complexity
- ultra-low cost
- Long battery life



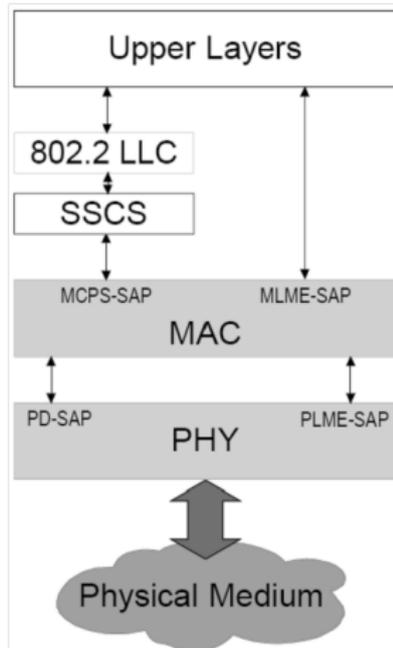
# Protocol Architecture

## PHY layer spec.

- 16 channels in 2.45 GHz ISM band
- max. bitrate: 250 kb/s
- max. symbol rate: 62.5 ksymbol/s

## MAC layer spec.

- CSMA-CA channel access
- Optional allocation of guaranteed time slots (GTSs)
- Beaconing for sync.
- Duty-cycled superframe structure
- MAC layer security



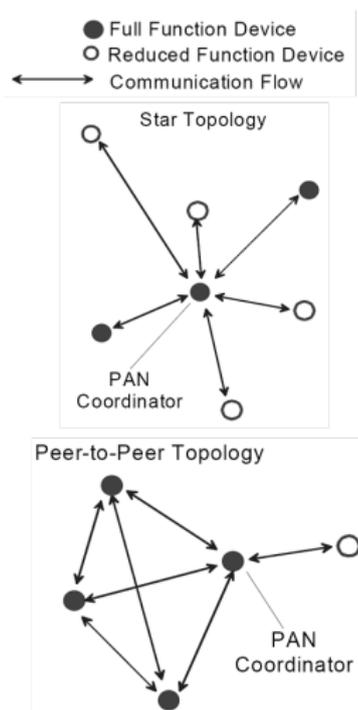
# Devices and Topologies

## Two device types

- Full-function device (FFD)
  - PAN coordinator, coordinator, or device
- Reduced-function device (RFD)
  - minimal resources
  - talk only to an FFD

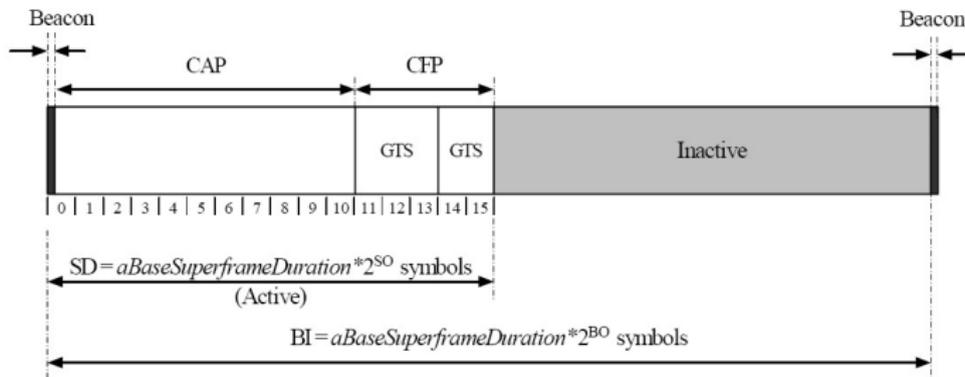
## Two topologies

- Star topology
- Peer-to-peer topology (mesh)
  - Cluster-tree



# Superframe Structure

- Periodical beacons sent by coordinators for sync.
- Active portion: channel access
  - Contention Access Period (CAP):
    - slotted CSMA-CA
  - Contention Free Period (CFP):
    - TDMA-like Guaranteed Time Slot (GTS)
- Inactive portion: go to sleep

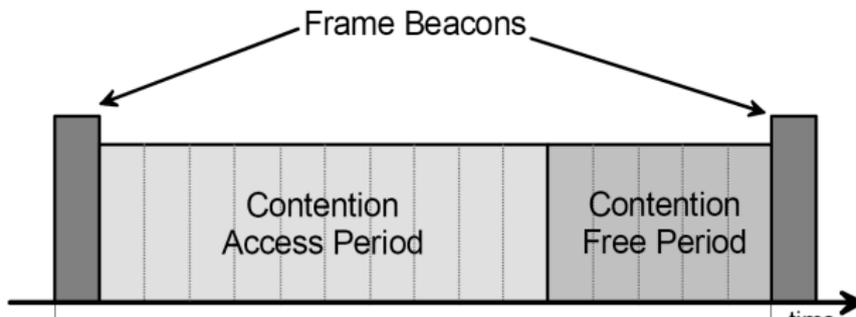


# Standard Protocol Behavior

Only GTS can provide guaranteed latency boundary

## Limitations in IEEE 802.15.4 GTS for studied scenario

- Maximum seven GTSs allowed – we need e.g. 20
- A minimum CAP length of 440 symbols required – introduce extra latency (7.04 ms/BI)
- One GTS must consist of an integer number of contiguous superframe slots – wast of bandwidth when transmitting e.g. 1 byte payload



# Removal of Limitations

- Only GTSs in the superframe, no CAP and inactive period  
 $\Rightarrow I_{CAP} = I_{SLP} = 0$
- Each GTS allocated with an exact bandwidth for one complete transaction  $\Rightarrow I_{GTS} = I_{TR} = I_D + I_{SIFS}$

$$I_{BI} = I_B + I_{SIFS} + I_{CAP} + n \times I_{GTS} + I_{SLP}$$

$$d_{BI} = I_{BI} / (62.5 \text{ ksymbols/s}) = 17.376 \text{ ms (for 20 nodes)}$$

Table: Duration Parameters

Symbol	Description	Value
$I_B$	length of beacon transmission	34 symbols
$I_D$	length of data transmission	40 symbols
$I_{SIFS}$	short interframe space	12 symbols
$I_{TR}$	length of one transaction	52 symbols

# Worst Case Estimation

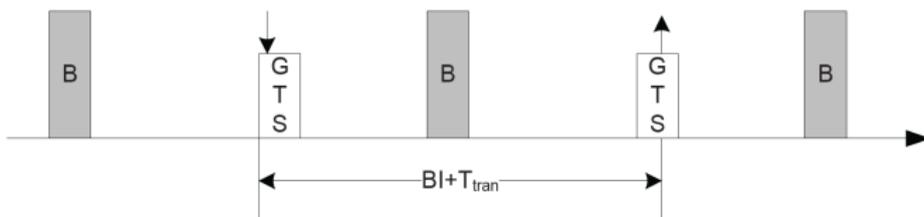
As usual means to calculate delay boundary

## Assumption

Each device possesses its own GTS slot in the superframe and keeps tracking the beacon

## Worst case in the studied scenario

The maximum latency occurs when a message is generated by a device during its own GTS slot slightly after the time point, at which it can be transmitted in the current GTS, and has to wait for an extra beacon interval.



# Analysis Results

## Guaranteed latency boundary

$$I_{GUA} = I_{BI} + I_{TR} = I_B + (n + 1) \times I_{SIFS} + n \times I_D$$

$$d_{GUA} = I_{GUA} / (62.5 \text{ ksymbols/s})$$

## In the case of studied scenario

$$d_{BI} = 17.376 \text{ ms}, d_{TR} = 0.832 \text{ ms}$$

$$d_{GUA} = 18.208 \text{ ms} > 10 \text{ ms}$$

## Analysis

$I_D$  and  $I_{SIFS}$  dominate the boundary value

- $I_{SIFS} = 12 \text{ symbols}$
- $I_D = 40 \text{ symbols} = 20 \text{ bytes}$   
 ⇒ Transmission of one byte payload needs 38 symbols (19 bytes) overheads added by MAC and PHY (too large)

# Overview

## Design goals

- Guaranteed latency upper-boundary
  - 20 nodes, 1 byte payload, max. 10 ms latency
- Energy-efficiency consideration
- Hardware compatibility

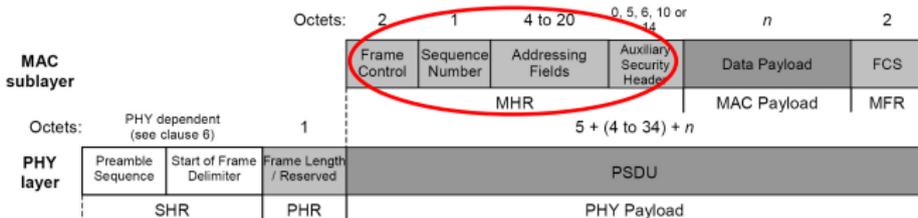
## General features

- Original IEEE 802.15.4 PHY
- TDMA-based superframe structure
- Largely reduced MAC overheads



# Data Frame w/o MAC Header

- Originally 19 bytes overheads added to one byte payload



## Abandon MAC header

- Pre-allocation -> implicit addressing
- No direct ACK -> SN unneeded
- Other fields dispensable

## Group Acknowledgment

PAN coordinator sends a group ACK in next beacon

# Parameters Recalculation

For studied scenario: 20 nodes, 1 byte payload

**Table:** Reduced Duration Parameters

Symbol	Standard w/o limitations	Improved protocols
$I_D$	40 symbols	18 symbols
$I_{SIFS}$	12 symbols	4 symbols
$I_{TR}$	52 symbols	22 symbols
$I_{BI}$	1246 symbols	494 symbols

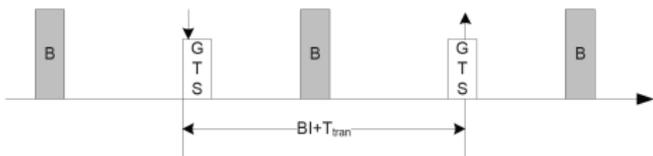
## Beacon interval

$$I_{BI} = I_B + 2 \times I_{IFS} + n \times I_{GTS} + (n - 1) \times I_{SIFS}$$

$$d_{BI} = (I_B + 2 \times I_{IFS} + 20 \times I_D + (20 - 1) \times I_{SIFS}) / (62.5 \text{ ksymbols/s})$$

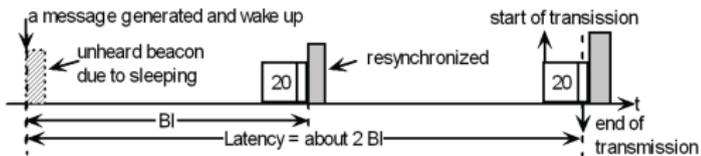
# Performance Analysis

- Beacon tracking enabled



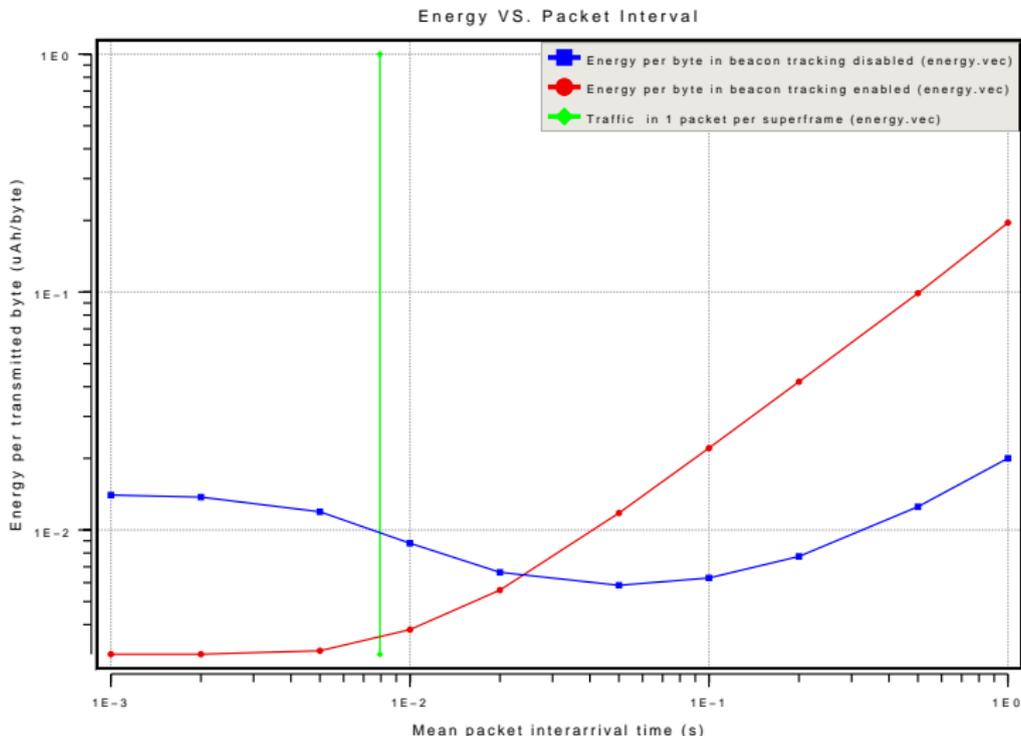
Guaranteed max latency:  $BI + t_{TRAN} = 8.3 \text{ ms}$

- Beacon tracking disabled



Guaranteed max latency:  $2BI = 15.81 \text{ ms}$

# Simulation Results for Energy Consumption

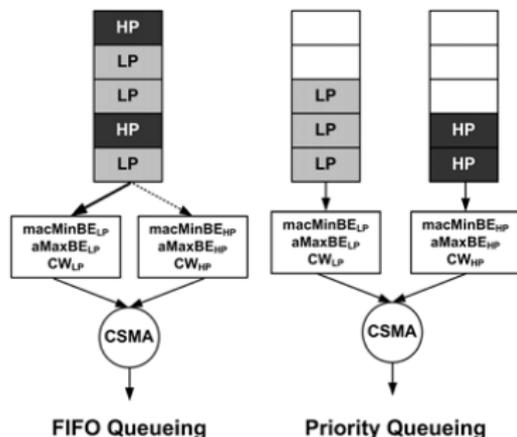


# Low-latency IEEE 802.15.4/CSMA

## Priority-based scheme providing differentiated QoS

- Priority queue for two types of data packets
  - High priority (HP) packets
  - Low priority (LP) packets
- Differentiated CSMA/CA parameters
  - One-time CCA VS. two-times CCA
  - $CW_{HP} \leq CW_{LP}$
  - $macMinBE_{HP} \leq macMinBE_{LP}$
- Priority toning

Slotted CSMA/CA with Service Differentiation



# Thank you and questions?